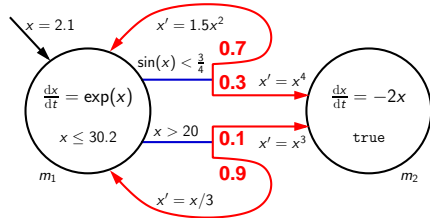


## Analysis of Probabilistic Hybrid Automata



Probability of  $\Diamond(m_2 \wedge x \leq -9.8)$ ?

## Stochastic SMT-based Model Checking

$\exists trans \in \{1, 2\} : \mathcal{V}_{[1 \rightarrow 0.7, 2 \rightarrow 0.3]} prob_1 \in \{1, 2\} :$

$\mathcal{V}_{[1 \rightarrow 0.1, 2 \rightarrow 0.9]} prob_1 \in \{1, 2\} : \dots$

$\wedge \left( (m_1 \wedge trans = 1 \wedge \sin(x) < \frac{3}{4} \wedge prob_1 = 1) \right. \\ \left. \implies (x' = 1.5x^2 \wedge m'_1) \right)$

$\wedge \left( (m_1 \wedge trans = 1 \wedge \sin(x) < \frac{3}{4} \wedge prob_1 = 2) \right. \\ \left. \implies (x' = x^4 \wedge m'_2) \right)$

$\wedge \dots$

### State of affairs

### Future work

Discrete-time /  
Scheduled event

Continuous-time

Bounded MC

Full MC  
(stochastic interpolation)

Reachability

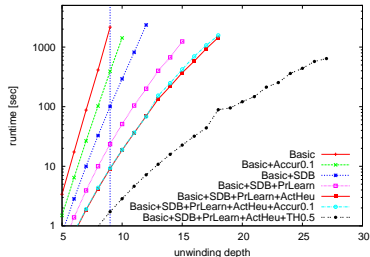
Expressive logics  
Probabilistic stability

Probability results

Counter-examples

### SSMT Algorithm

- Traversing quantifier tree
- SMT solver for quantifier-free subproblems
- Aggressive pruning rules for efficiency



# Quantitative Program Semantics and Analysis

- Probabilistic Programming **Languages**
  - pccp [AGP97, ICCL98, MFCS98]
  - pKLAIM [Coordination04, SecCo04]
  - pCHAM [FMCO05]
  - pLambda [JLC05]
  - pWhile [APLAS07, ICICS08]
- Denotational and Operational **Semantics**
- Quantitative Static **Program Analysis**
- Quantitative Aspects in **Computer Security**
- **Implementation** Based on Linear Algebra

# Quantitative Program Semantics and Analysis

- Probabilistic Programming **Languages**
- Denotational and Operational **Semantics**
  - Probabilities and Non-Determinism
  - Discrete Time vs Continuous Time
  - Operator Algebras [MFCS98,MFCSIT04]
  - Compositional Semantics [APLAS07]
- Quantitative Static **Program Analysis**
- Quantitative Aspects in **Computer Security**
- **Implementation** Based on Linear Algebra

# Quantitative Program Semantics and Analysis

- Probabilistic Programming **Languages**
- Denotational and Operational **Semantics**
- Quantitative Static **Program Analysis**
  - Probabilistic Abstract Interpretation
  - Moore-Penrose Pseudo Inverse [PPDP00,LNCS4444]
  - Syntax Directed Semantics [JFP05,APLAS08]
- Quantitative Aspects in **Computer Security**
- **Implementation** Based on Linear Algebra

# Quantitative Program Semantics and Analysis

- Probabilistic Programming **Languages**
- Denotational and Operational **Semantics**
- Quantitative Static **Program Analysis**
- Quantitative Aspects in **Computer Security**
  - Approximate Confinement [AGP00,CSFW02,CONCUR03]
  - Hypothesis Testing [CSFW02,TCS05,JCS04]
  - Most Effective Attacker [CSFW02,SAS02]
  - Timing Attacks and (Counter)Measures [SAS02,ICICS08]
  - Probabilistic Program Transformation [JLAP07,ICICS08]
- **Implementation** Based on Linear Algebra

# Quantitative Program Semantics and Analysis

- Probabilistic Programming Languages
- Denotational and Operational Semantics
- Quantitative Static Program Analysis
- Quantitative Aspects in Computer Security
- Implementation Based on Linear Algebra
  - Sparse Matrices
  - Tensor Product
  - Octave and OCaml

## Research interests

- ▶ process algebra for hybrid systems: HYPE
  - ▶ discrete and continuous behaviour
  - ▶ permits modelling of individual flows
  - ▶ compositionality as an important feature
  - ▶ Galpin, Hillston & Bortolussi, MFPS 2008, ENTCS 218, 2008
- ▶ spatial stochastic process algebra
  - ▶ addition of spatial aspects to PEPA
  - ▶ physically distributed systems, computer and biological
  - ▶ Galpin, AINA 2009, to appear
- ▶ semantic equivalences in discretised systems
  - ▶ behavioural equivalence between two discrete models of the same system

## Application of logic to control problems using Multi-dimensional System co-Engineering

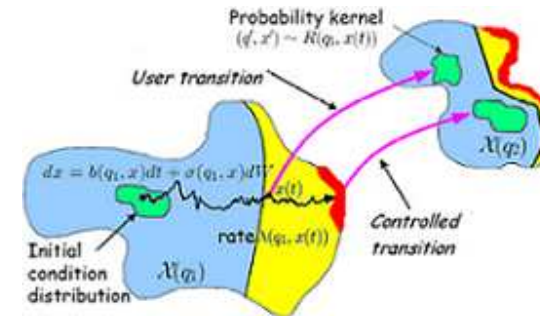
Manuela L. Bujorianu and Marius C. Bujorianu

Centre for Interdisciplinary Computational and Dynamical Analysis, University of Manchester

*Multi-dimensional system co-engineering* [1], abbreviated **MscE** [2], is a modeling framework that combines formal, mathematical and control engineering. It contains

- a **reference model**, called *colored stochastic hybrid systems (cSHS)*,
- **specification logics** like **SafAL** (the Safety analysis logic), **Hil** (the Hilbertean logic), **CSL** (extended for continuous processes), united in the paradigm of *Hilbertian Formal Methods* [5]) for specification of safety and reachability properties, and
- a toolset for **formal and dependable verification**.

The cSHS combines a very general model of *stochastic hybrid systems (SHS)* [3] analysis information, modeled as colors. A SHS describes the evolution of a *hybrid* under the influence of stochastic perturbations. A HS consists of a digital controller that can evolve in different modes, modeled as (deterministic or stochastic) continuous dynamical system. Moreover, the cSHS model is extended to communicating autonomous multi-agent systems [6]. Instances of this framework have been used to model and analyze systems forms air traffic control [4], [7]. Currently we explore the issues of modeling, control, coordination and verification for aerospace systems, such as (formations of autonomous) satellites.



with runtime  
system (HS)  
and a plant

[1] Multi-dimensional System Co-Engineering  
<http://personalpages.manchester.ac.uk/staff/Manuela.Bujorianu/MScE.htm>

[2] M.C. Bujorianu, M.L. Bujorianu and H. Barringer “A Formal framework for user centric control of multi-agent cyber-physical systems” in Michael Fisher, Fariba Sadri, Michael Thielscher Proceedings of the 9th International Workshop on Computational Logic in Multi-Agent Systems (CLIMA), Springer Verlag LNCS, 2009

[3] M.L. Bujorianu. “Extended Stochastic Hybrid Systems and their Reachability Problem” In: Proceedings Hybrid Systems: Computation and Control, HSCC 2004, pp. 234-249, Springer LNCS vol. 2993, 2004,

[4] Giordano. Pola, Manuela L. Bujorianu, John Lygeros and Maria Di Benedetto “Stochastic Hybrid Models: An Overview with Application to

Air Traffic Management” In: 1st IFAC Conf. on Analysis and Design of Hybrid Systems, ADHS 2003, pp. 45-50, 2003

[5] M.C. Bujorianu and M.L. Bujorianu “Towards Hilbertian Formal Methods” Proceedings of Application of Concurrency to System Design ACSD’07, IEEE Computer Society Press, pp. 240-241, 2007

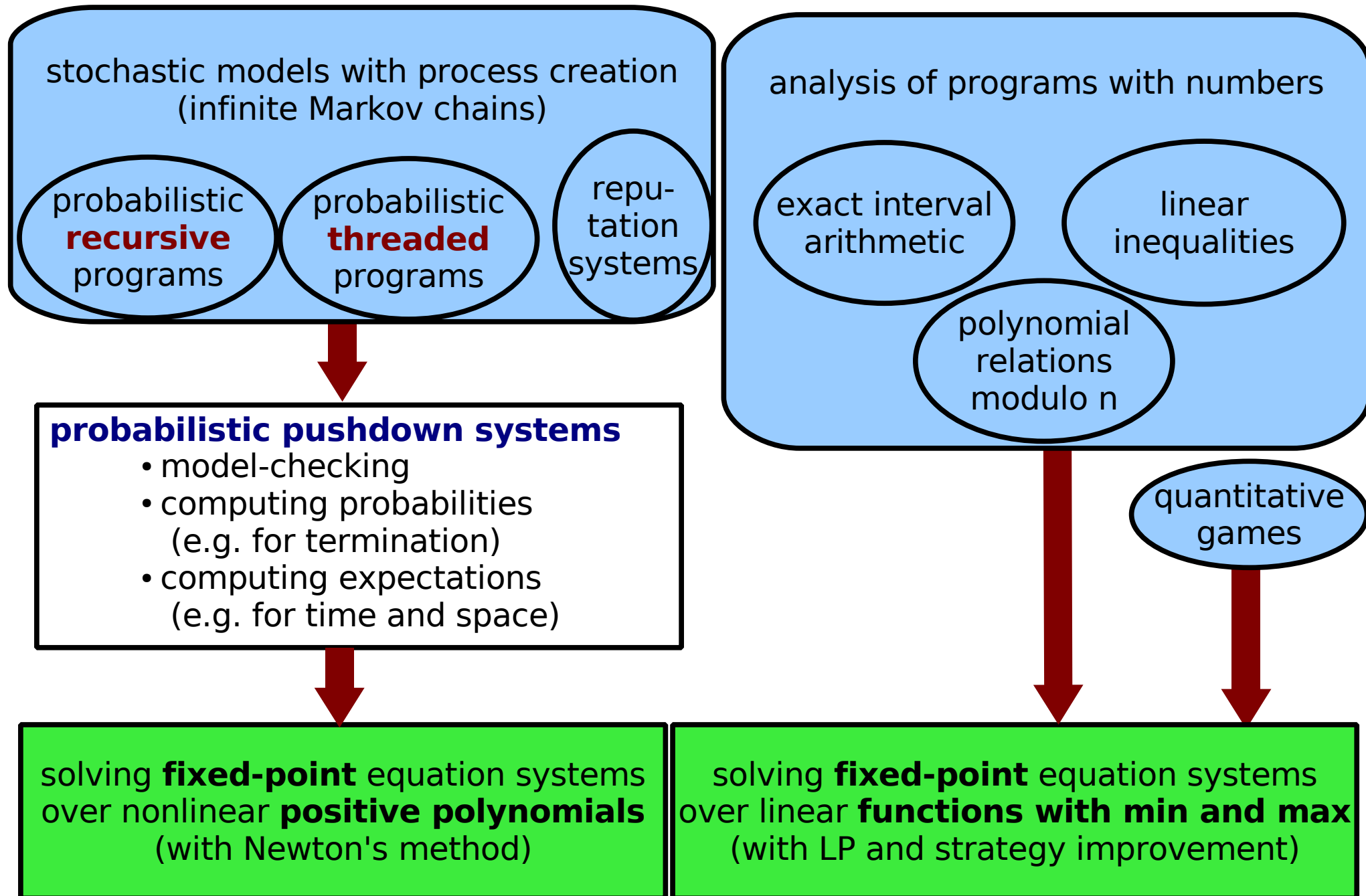
[6] M.C. Bujorianu and M.L. Bujorianu, and Savi Maharaj “Distributed Stochastic Hybrid Systems” In Horacek, P., Simandl, M. and Zitek, P., Proceedings of IFAC 2005, Elsevier Science Press 2005

[7] Hybridge Distributed Control and Stochastic Analysis of Hybrid Systems Supporting Safety Critical Real-Time Systems Design  
<http://www2.nlr.nl/public/hosted-sites/hybridge/>



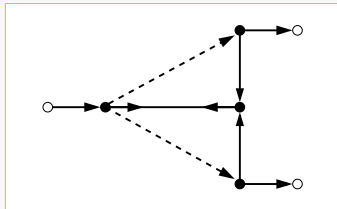
# Quantitative Analysis at TU München

(Groups of Javier Esparza and Helmut Seidl)



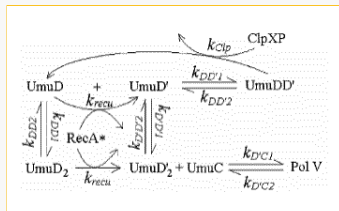
# MLQA building blocks from CWI and TU/e

## Stochastic Reo



- SOA and QoS
- Connector synthesis
- Dynamic adaptor modification

## Formal cell processes



- Systems biology
- Stochastic process languages
- ODE vs. Gillespie vs. PCTL

stream calculus and coalgebra & control theory as discussed by Milad Niqui

# Coalgebras

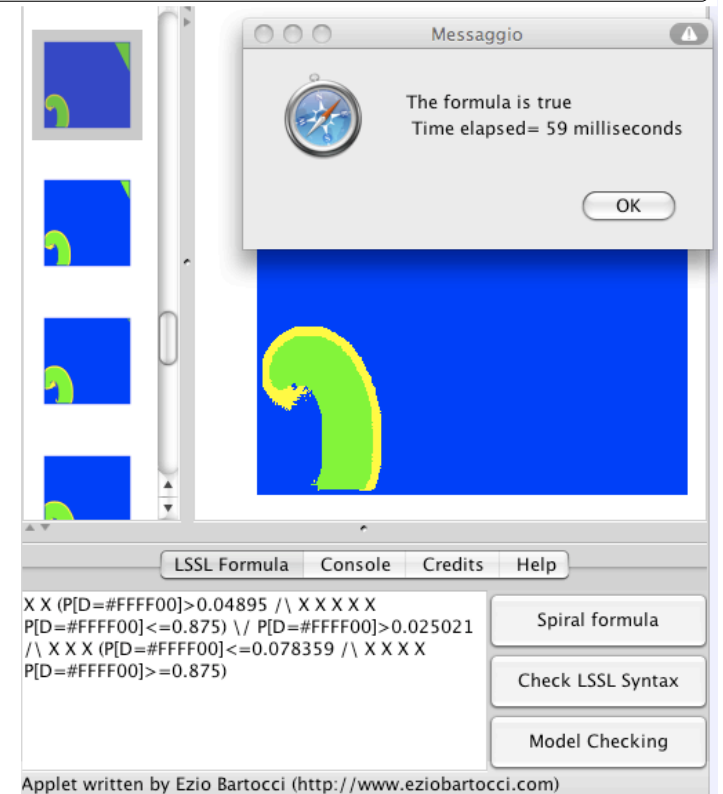
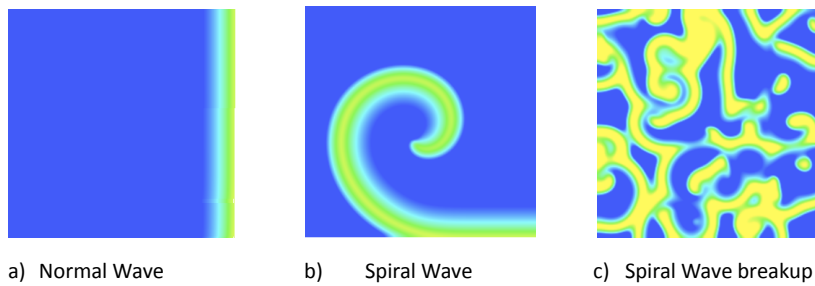
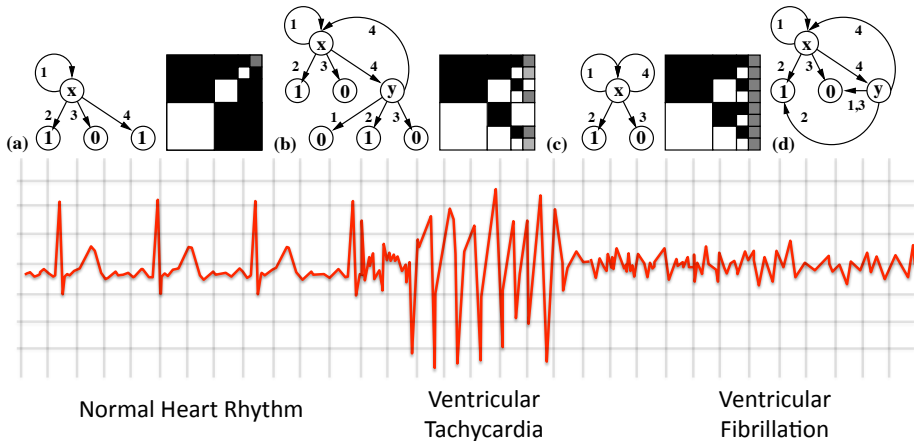
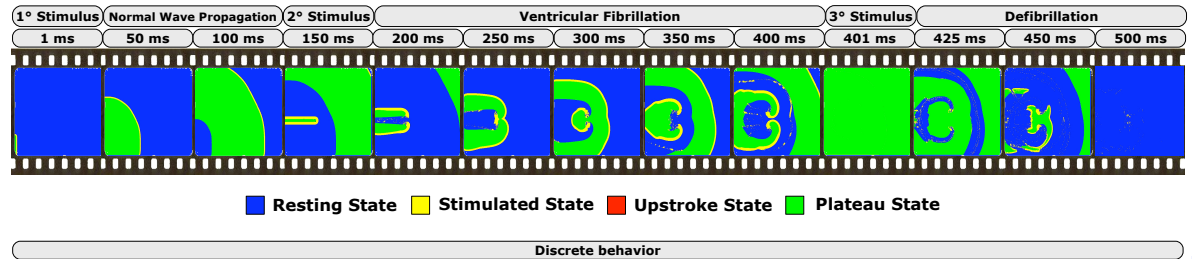
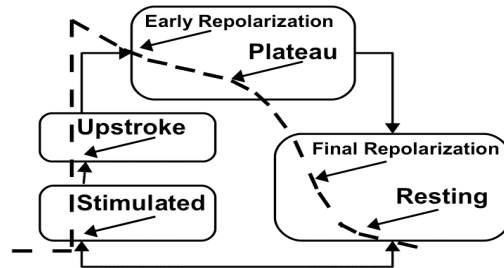
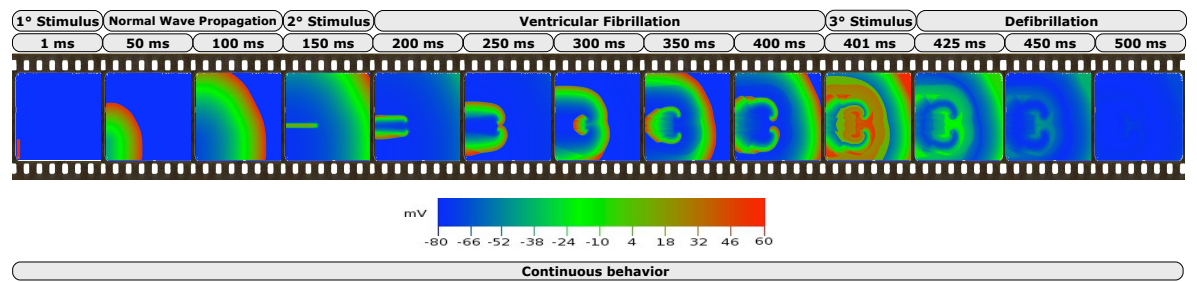
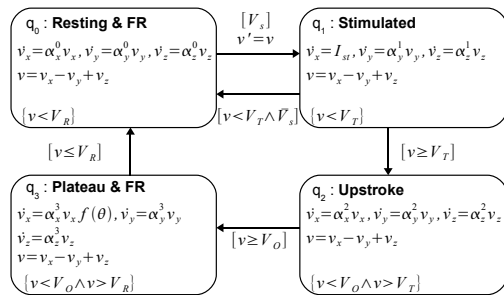
- Realisation of formal power series
- Stream calculus
- Exact arithmetic
- Coinduction
- (Bi)simulation & trace
- Temporal logic
- Automata

## Embedded Systems

- Linear & rational systems
- Discrete event systems
- Continuous-time systems
- Smooth systems
- Hybrid systems

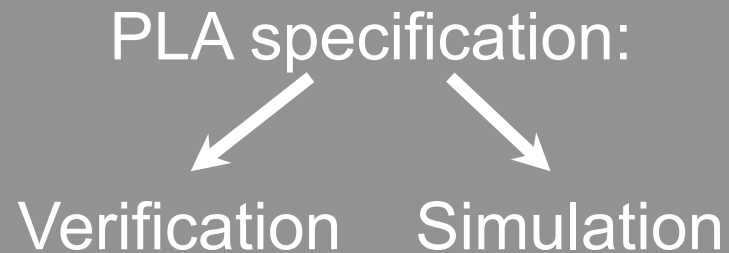
## Service Oriented IT

- Composition & hiding
- Coordination from outside
- Coinductive behaviour types
- Distributed systems





PLA model = <Piece-linear Markov process,  
aggregate system, controlling sequences>



Software tools for PLA:

- Verifier;
- Simulator;
- Automated model creator for Markov processes.

Applications:

- telecommunication protocols;
- business systems;
- biomedicine.

Modifications:

- hybridPLA;
- MarkovPLA;
- dynPLA.

## RESEARCH INTERESTS

- hybrid process calculi in the context of

## FLYING SENSORS

(Present)

- ▷  $n > 1$  picosatellites fulfill tasks using cooperative and dynamic trajectory coordination
- ▷ *some problems*: energy, complex algorithms, swarm behaviour
- ▷ to be sponsored by the DFG?

- stochastic process calculi in the context of

## ANALYSIS OF PEER-TO-PEER ALGORITHMS

(Past)

- ▷ modeling using a distributed stochastic process calculus
- ▷ specification using the markov chain based logic CSL
- ▷ verification by model checking after statespace reductions

General process algebraic specification

$$X = a \cdot b \cdot X + c \cdot X$$

## General process algebraic specification

$$X = {}^1a \cdot {}^2b \cdot X + {}^1c \cdot X$$



## General process algebraic specification

$$X = {}^1a \cdot {}^2b \cdot X + {}^1c \cdot X$$

## Linear format

$$\begin{aligned} X(pc) = \\ & pc = 1 \Rightarrow a \cdot X(2) \\ & + pc = 1 \Rightarrow c \cdot X(1) \\ & + pc = 2 \Rightarrow b \cdot X(1) \end{aligned}$$

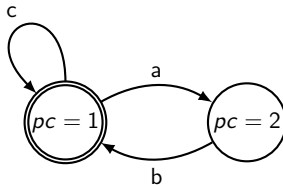
# Symbolic Translation of Stochastic Processes

## General process algebraic specification

$$X = {}^1a \cdot {}^2b \cdot X + {}^1c \cdot X$$

## Linear format

$$\begin{aligned} X(pc) = \\ & pc = 1 \Rightarrow a \cdot X(2) \\ & + pc = 1 \Rightarrow c \cdot X(1) \\ & + pc = 2 \Rightarrow b \cdot X(1) \end{aligned}$$



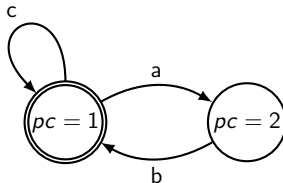
# Symbolic Translation of Stochastic Processes

## General process algebraic specification

$$X = {}^1a \cdot {}^2b \cdot X + {}^1c \cdot (0.5 : d \cdot X, 0.5 : e \cdot X)$$

## Linear format

$$\begin{aligned} X(pc) = \\ & pc = 1 \Rightarrow a \cdot X(2) \\ & + pc = 1 \Rightarrow c \cdot X(1) \\ & + pc = 2 \Rightarrow b \cdot X(1) \end{aligned}$$



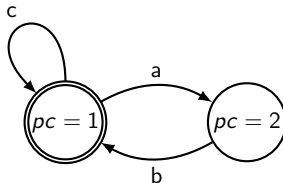
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## Linear format

$$\begin{aligned} X(pc) = \\ & pc = 1 \Rightarrow a \cdot X(2) \\ & + pc = 1 \Rightarrow c \cdot (0.5 : X(3), 0.5 : X(4)) \\ & + pc = 2 \Rightarrow b \cdot X(1) \end{aligned}$$



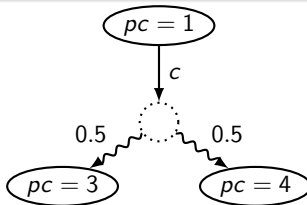
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# To use Model checking for Control design

- Problem:
- Hybrid systems are **undecidable**
  - Decidable models of CS are currently **unusable**

MLQA mission statement:

“This spans [...] *resource usage* (e.g. ‘the control system rotates and adjusts the windmill such that at least 60 % of the potential wind energy is utilised’).”

- Highly **non-linear** system !
- Non-linear hybrid systems are undecidable
- State of the art: *Abstraction* to **discrete** system
- or to discrete-time or real-time system (e.g. **timed automata**)
- Problem: abstraction **too coarse**  $\Rightarrow$  unusable for most examples
- Badly need **finer abstractions**
- Good candidate: **Priced timed automata**
- Algorithms; tool support

Watch us in Aalborg...