

Aim : modelling and estimating resource usage of software



Generic model for resources



Quantitative semantics

$q(i, a_2, b_2) = 12 \otimes 8$
Accumulation of costs along a path
 $q(a_1, b_1) = 2 \oplus 3$
Maximum cost between two paths
 $\sqrt[3]{q(i, a_1, b_1, c1)} = \sqrt[3]{8 \otimes (2 \oplus 3) \otimes 4}$
Average cost of a path
 $\mathbb{1}$: neutral cost
 \emptyset : impossibility of a transition

Labeled transition system

$$\rho(M) = \bigoplus_{k=1}^{|\Sigma|} \sqrt[k]{tr M^k}$$

Long-run cost
maximum average along traces

$$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 8 & 0 & 0 & 5 & 0 & 0 & 0 \\ 0 & 3 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 4 & 0 & 0 & 0 & 0 \\ 12 & 0 & 0 & 0 & 0 & 0 & 6 \\ 0 & 0 & 0 & 0 & 8 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$

Transition system as a matrix M :
linear operator semantics
Hint : see \oplus and \otimes as max and + over reals

Abstraction

Abstract transition system

$$\begin{pmatrix} \mathbb{1} & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & \mathbb{1} & 0 & 0 & 0 & \mathbb{1} & 0 \\ 0 & 0 & \mathbb{1} & 0 & 0 & 0 & \mathbb{1} \\ 0 & 0 & 0 & \mathbb{1} & 0 & 0 & \mathbb{1} \end{pmatrix}$$

Abstraction matrix α

$$\begin{pmatrix} 0 & 0 & 0 & 0 \\ 12 & 0 & 0 & 6 \\ 0 & 8 & 0 & 0 \\ 0 & 0 & 4 & 0 \end{pmatrix}$$

Abstract semantics $M^\#$

The concrete semantics matrix M is abstracted into a smaller one $M^\#$ thanks to an abstraction operator α that merges states

Correctness condition for an abstraction
 $\alpha \circ M \leq M^\# \circ \alpha$

Derivation of the « best » abstract semantics :
 $M^\# = \alpha \circ M \circ \alpha^{-1}$

Soundness of abstractions

Theorem : $\rho(M) \leq \rho(M^\#)$

Correct abstractions safely over-approximate long-run cost computations

Example of a set lattice and its associated powerset

Improvements in abstraction design

- reconciling classical and quantitative linear models
- reusing well-known abstractions (intervals, polyhedra ...)
- vector representation that takes two order relations into account

Classical abstract interpretation

- concrete and abstract semantics are linked via a Galois connection
- partial order on states

$pc \text{ setmode } m' \text{ } pc'$
 $(pc, \rho, m) \rightarrow^{c_m(m, m')} (pc', \rho, m')$

Quantitative semantics rule

```

setsize b n;
setsize tmp n;
if n < 100 then setmode A else setmode D
apply copy_to b, tmp;
    
```

Explicit energy modes management

Application : computing in various energy modes

Energy mode	Time (in cycles)	Cost (wrt array size)
A	5	0.2
C	2	$\lambda n.n$
D	1	$\lambda n.3. * n$
E	0	$\lambda n.5. * n$

Cost of an array operation

Energy mode		Threshold (array size)		
High	Low	100	500	900
A	B	3.6	3.6	3.6
A	C	13.7	50.9	88.1
A	D	39.9	191.6	343.3
A	E	81.2	405.1	728.9
B	C	13.7	50.9	88.1
B	D	39.9	191.6	343.3
B	E	81.2	405.1	728.9
C	D	97.5	191.6	343.3
C	E	97.5	405.1	728.9
D	E	381.6	405.1	728.9

Example of experimental results